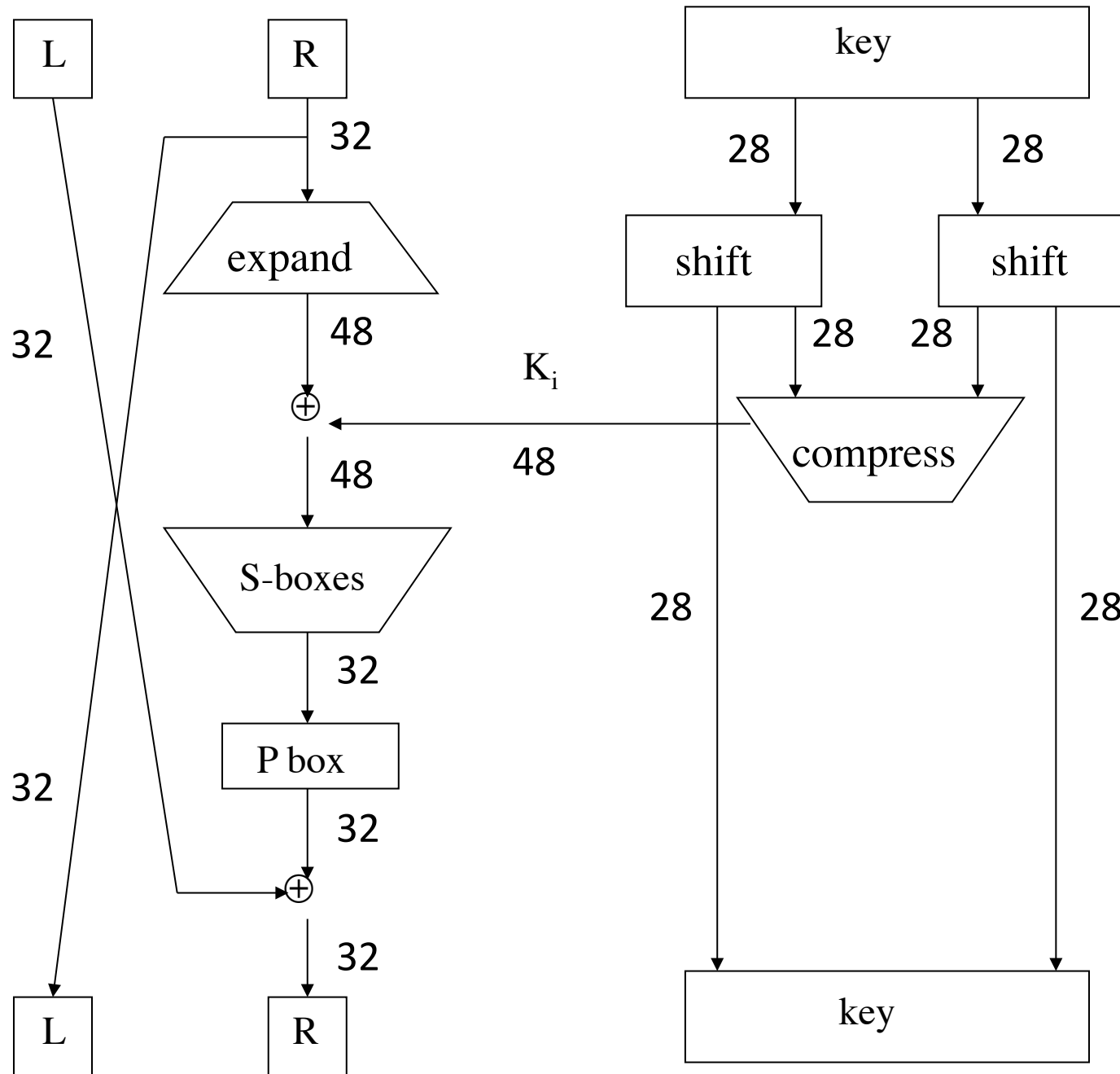


Triple DES

More on DES

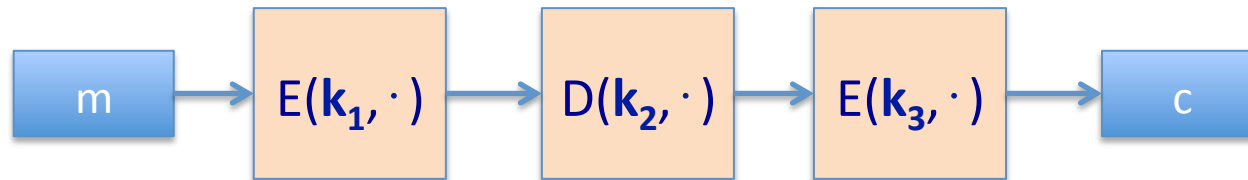
- DES Review
- Improve DES
 - 2DES
 - 3DES wrong
 - 3DES 3key
 - 3DES 2key



One
Round
of
DES

Triple DES

- **Triple DES** or **3DES** with three keys
 - $C = E(D(E(P, K_1), K_2), K_3)$
 - $P = D(E(D(C, K_3), K_2), K_1)$



- Keying options
 - option 1: all different (strongest)
 - option 2: $K_1 = K_3$ (most common)
 - option 3: $K_1 = K_2 = K_3$ (=DES)

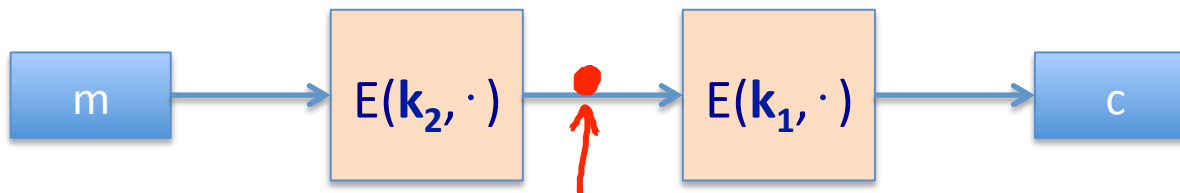
Triple DES

- Q1. Why not double DES (2DES)?
 - Meet In The Middle (MITM) Attack
 - takes 2^{63}
- Q2. Why not $E \rightarrow E \rightarrow E$ instead $E \rightarrow D \rightarrow E$?
 - Backward Compatibility (3DES option 3 = DES)
- Q3. MITM against 3DES?
 - Possible, but takes 2^{118}

Why not double DES?

- Define $2E((k_1, k_2), m) = E(k_1, E(k_2, m))$

key-len = 112 bits for DES



Find (k_1, k_2) s.t.
 $E(k_1, E(k_2, M)) = C$
 Equivalently:
 $E(k_2, M) = D(k_1, C)$

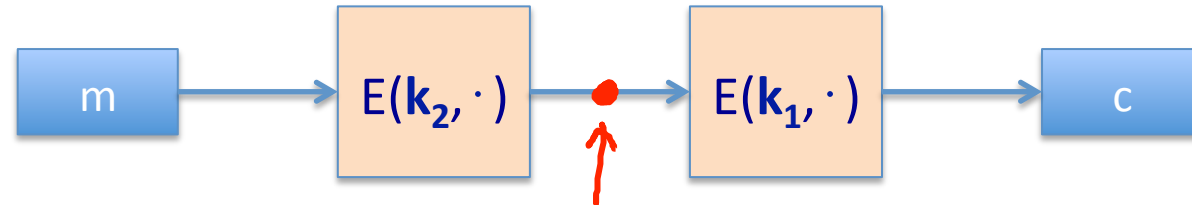
Attack: $M = (m_1, \dots, m_{10})$, $C = (c_1, \dots, c_{10})$

- step 1: build table.

sort on 2nd column

$k^0 = 00\dots00$	$E(k^0, M)$	} 2^{56} entries
$k^1 = 00\dots01$	$E(k^1, M)$	
$k^2 = 00\dots10$	$E(k^2, M)$	
$\vdots$	$\vdots$	
$k^N = 11\dots11$	$E(k^N, M)$	

Meet in the middle attack



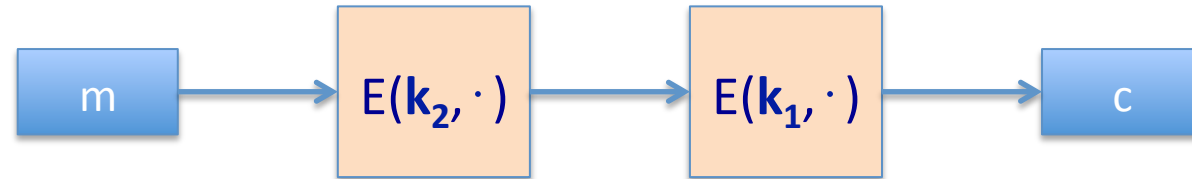
Attack: $M = (m_1, \dots, m_{10})$, $C = (c_1, \dots, c_{10})$

$k^0 = 00\dots00$	$E(k^0, M)$
$k^1 = 00\dots01$	$E(k^1, M)$
$k^2 = 00\dots10$	$E(k^2, M)$
$\vdots$	$\vdots$
$k^N = 11\dots11$	$E(k^N, M)$

- step 1: build table.
- Step 2: for all $k \in \{0,1\}^{56}$ do:
test if $D(k, C)$ is in 2nd column.

if so then $E(k^i, M) = D(k, C) \Rightarrow (k^i, k) = (k_2, k_1)$

Meet in the middle attack



Time = $2^{56}\log(2^{56}) + 2^{56}\log(2^{56}) < 2^{63} \ll 2^{112}$,
space $\approx 2^{56}$

Same attack on 3DES: Time = 2^{118} , space $\approx 2^{56}$

